

discrete log problem

give: P and P_A , find $a = \log_p P_A$,
 i.e. $P_A = aP$

naive method: try $P, 2P, 3P, \dots$ ↖ neutral element

assume that P has order N , i.e. $NP = 0P$

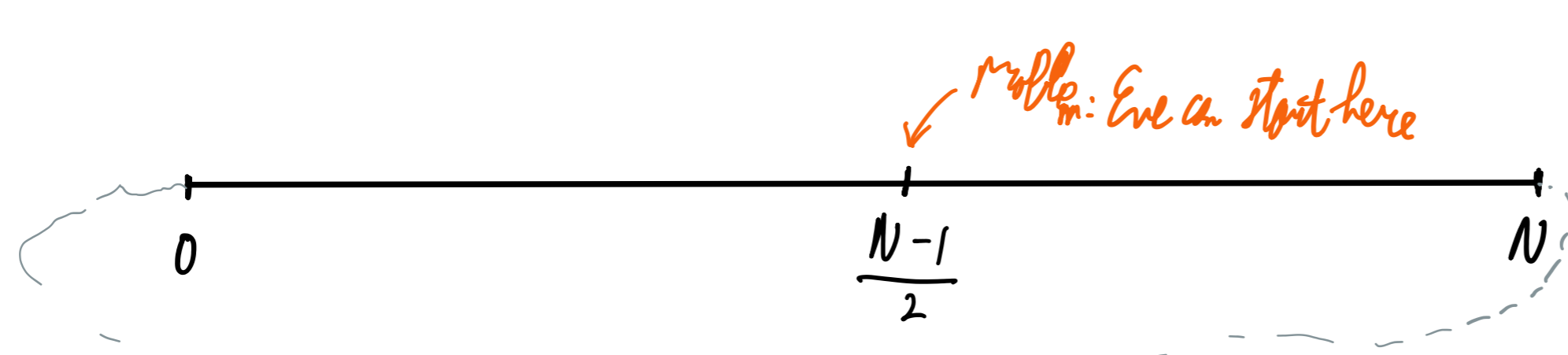
naive method takes at most $N-1$ additions

Can use $-P = (N-1)P$ ↗ (i.e. wrapping around) symmetrically to succeed in at most $\lceil \frac{N-1}{2} \rceil$ steps

If Eve is known to try $P, 2P, 3P,$

then a large a is safer.

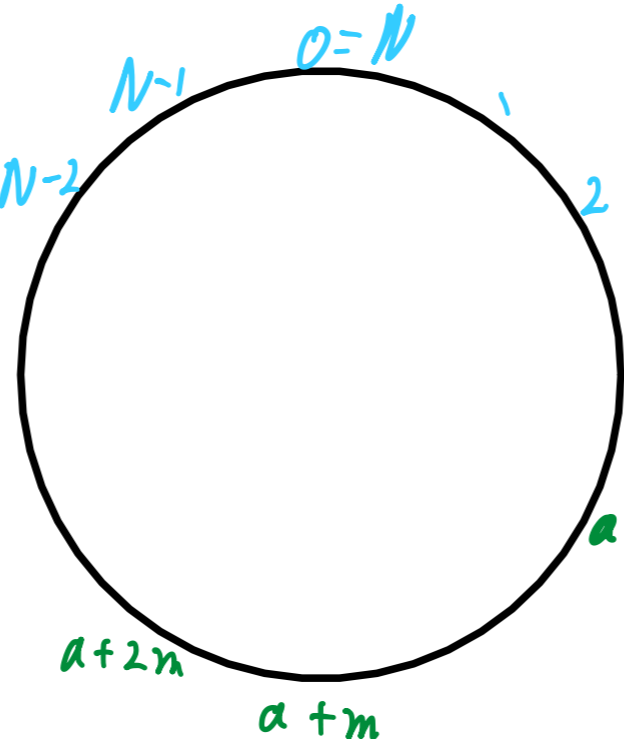
(as long as $a < \frac{N-1}{2}$ if Eve uses symmetry)



Eve can randomize P_A to avoid such "protections"

$P_A + rP$ for random r moves target DLP to $a+r \pmod N$

Eve can "blow up" the single target into 100 targets $P_A + r_1P, P_A + r_2P, \dots$ for random r_i .
 this improves her chance of finding one in $P, 2P, 3P$



nicer. Split interval $[0, N-1]$ into 100 equal pieces. let $m = \frac{N}{100}$, and pick targets $P_A, P_A + mP, P_A + 2mP, P_A + 3mP, \dots, P_A + 99mP$, so that one is in the first $\frac{N}{100} = m$ steps

(mp uses double-add method)

if $P_A + rP = cP$ for a c and r we know

(Eve computes $P, 2P, \dots$, at each step checks against all targets)

then $aP + rP = cP$, i.e.

$$a \equiv c - r \pmod N$$

The first DLP will be solved in $\frac{N}{100}$.

With 1000 targets, i.e. $m = \frac{N}{1000}$, solve in $\frac{N}{1000}$

With 10^i targets, $m = \frac{N}{10^i}$, solve in $\frac{N}{10^i}$

Warning: the first step takes 10^i computation and 10^i storage.
 the second step (compute $P, 2P, \dots, cP$ and compare to list) takes $\frac{N}{10^i}$ steps.

Do not choose i so large that $\frac{N}{10^i} < 10^i$. So, limit i so that $10^i \leq \sqrt{N}$

Storage is more expensive than computation, so limit the steps which incur storage.

This attack is (mostly, apart from swapping phase 1 and 2) the baby-step-giant-step attack.

Baby-step-giant-step attack (BSGS attack)

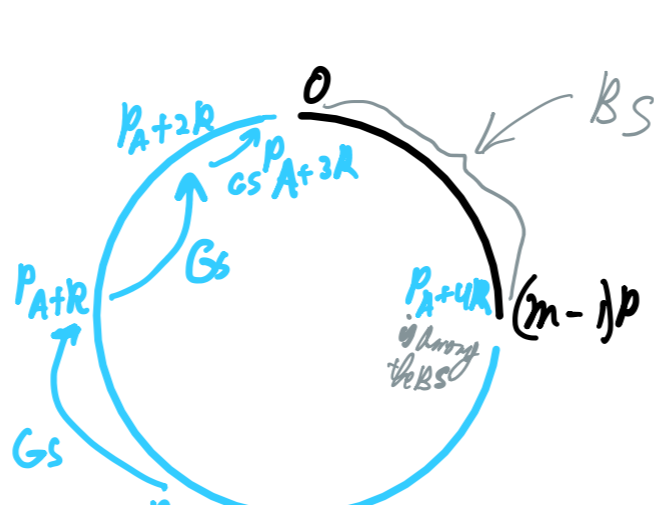
Give P, P_A , and N , $m = \lfloor \sqrt{N} \rfloor$

Baby steps: compute $0P, 1P, 2P, \dots, (m-1)P$ and store as (iP, i) indexed by iP .

Giant steps: compute $R = mP$.

$j=0$, $Q = P_A$
 while Q not in targets
 $j = j + 1$
 $Q = Q + R$

output j and i , where we have a match



$$P_A + 4R = iP$$

$$aP + 4mP = iP$$

$$a \equiv i - 4m \pmod N$$

↖ j ignored